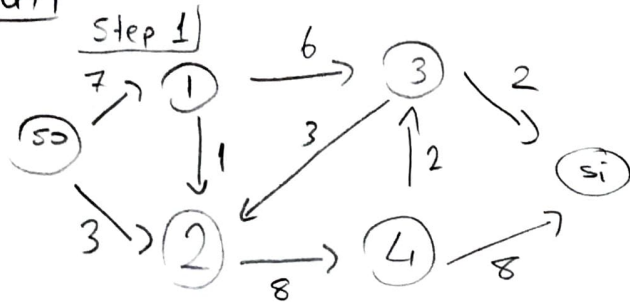
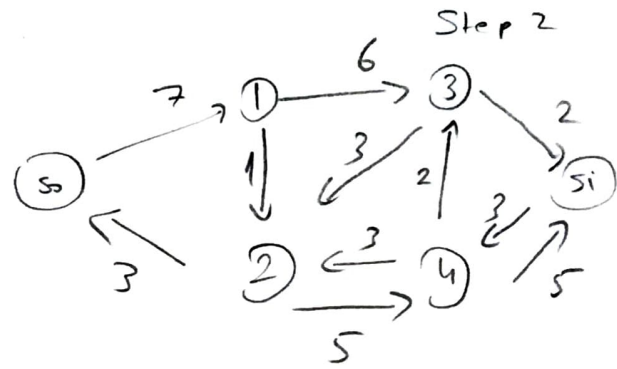


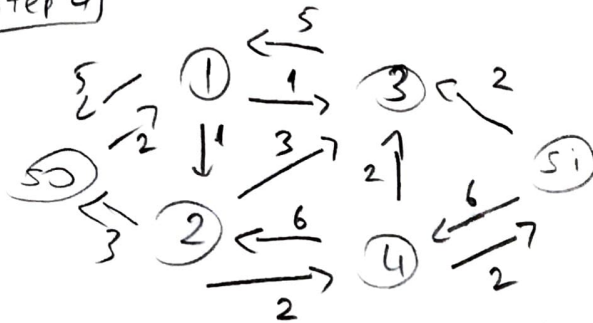
Q1)



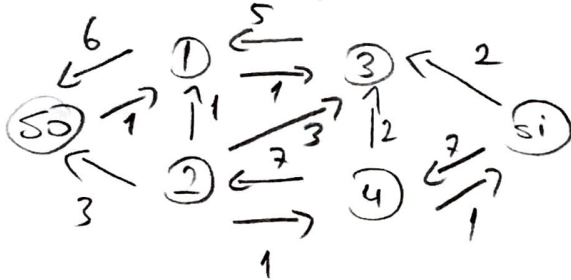
Send Flow of 3
through
 $s0 \rightarrow 2 \rightarrow 4 \rightarrow si$



Step 4)

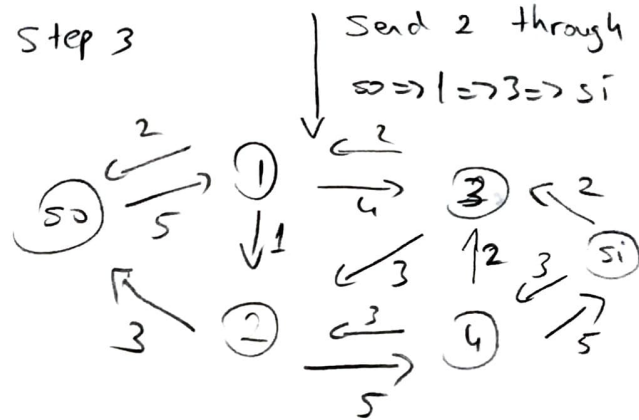


Flow of 1
through
 $s0 \rightarrow 1 \rightarrow 2 \rightarrow 4 \rightarrow si$



Flow of 3
 $s0 \rightarrow 1 \rightarrow 3$
 $\rightarrow 2 \rightarrow 4 \rightarrow si$

Step 3



(Augmenting path)
There is no directed path from $s0 \rightarrow si$. So this is our last step. (We looked for directed graphs and send the minimum over this path in each step)

Total flow is $3+2+3+1=9$

Therefore max flow is 9.

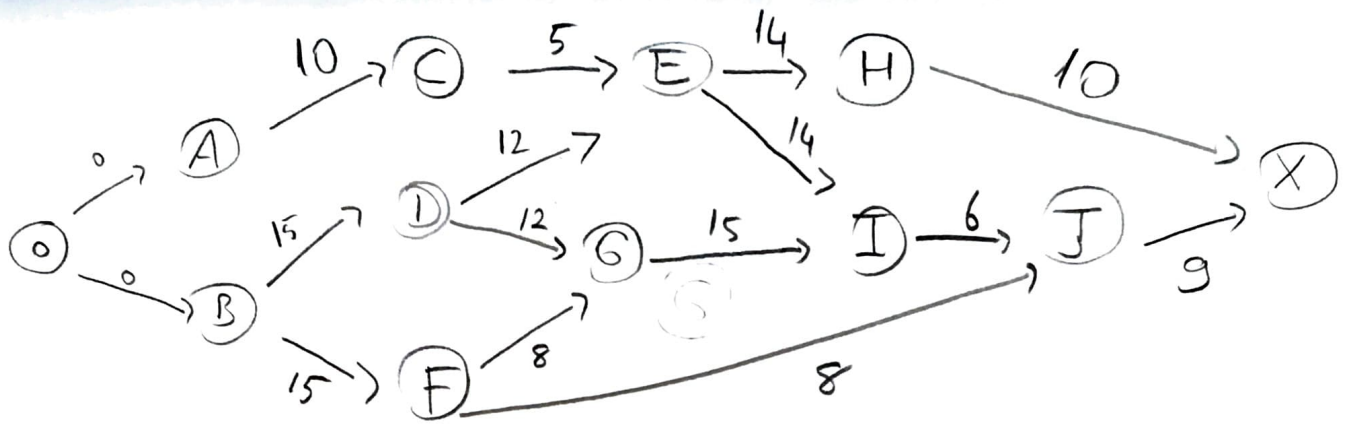
S-T cut is $S = \{s0, 1, 3\}$, $T = \{si, 2, 4\}$

(~~and~~ S = Nodes that we can go from $s0$, T is others?)

Capacity of this S-T cut is $(s0, 2) + (1, 2) + (3, 2) + (3, si) = 3 + 1 + 3 + 2 = 9$.

We cannot find a smaller S-T cut since max flow = min cut.

Q21



	ES	EF	LS	LF	Slack
A	0	10	13	23	13
B	0	15	0	15	0
C	10	20	23	28	13
D	15	27	15	27	0
E	27	41	28	42	1
F	15	30	19	27	4
G	27	39	27	42	0
H	41	51	47	57	6
I	42	57	42	57	0
J	48	54	48	57	0
X	57	66	57	66	0

ES of a point = max of its predecessors EF

EF = ES + activity time

LS = LF - activity time

LF = Min of LS of its later dependent points

Slack = LS - ES = LF - EF

Critical activities will be the ones that slack = 0

so they are B, D, G, I, J.

Duration is 57

e) Anything can be that has slack variable > 0. (A, C, E, F, H)

f) a_i = activity time activity i , x_i = start time of activity i

min x_x subject to $x_0 = 0$, $x_A \geq 0$, $x_B \geq 0$, $x_C \geq x_A + 10$, $x_D \geq x_B + 15$

$x_E \geq x_D + 12$, $x_E \geq x_C + 5$, $x_F \geq x_B + 15$, $x_G \geq x_D + 12$, $x_G \geq x_F + 8$, $x_H \geq x_E + 14$

$x_I \geq x_E + 14$, $x_I \geq x_G + 15$, $x_J \geq x_F + 8$, $x_J \geq x_I + 6$; $x_x \geq x_J + 9$, $x_x \geq x_H + 10$

~~$x_A, x_B, x_C, x_D, x_E, x_F, x_G, x_H, x_I, x_J, x_x \geq 0$~~

(set of arcs)

or shortly min x_x subject to $x_j \geq x_i + a_i$ for all $(i, j) \in \hat{A}$

$x_i \geq 0$ for all $i \in N$ $x_0 = 0$

(set of nodes)

Q3 | $f_k(d)$ = max reliability with components k and available d dollar.

r_{ij} : reliability of component i with j parallel units

c_{ij} : cost of using

$f_0(d) = 1$ for $d = 0, 1, \dots, 10$ and $f_i(0) = 0$ for $i = 1, 2, 3$.

$$f_i(d) = \max_{\substack{j=1,2,3 \\ c_{ij} \leq x}} \{ r_{ij} f_{i-1}(x - c_{ij}) \}$$

$$\text{So } f_1(1) = 0.6 ; f_1(2) = 0.8 ; f_1(3) = f_1(4) = \dots = f_1(10) = 0.9$$

$$f_2(1) = 0, f_1(1) = 0 ; f_2(2) = 0, f_1(2) = 0 ; f_2(3) = 0.7 \cdot f_1(0) = 0$$

$$f_2(4) = 0.7 \cdot 0.6 = 0.42 ; f_2(5) = \max \{ 0.8 f_1(0), 0.7 f_1(2) \} = 0.56$$

$$f_2(6) = \max \{ 0.8 f_1(1) = 0.48, 0.7 f_1(3) = 0.63 \} = 0.63$$

$$f_2(7) = \max \{ (0.9) f_1(1), 0.8 f_1(2), 0.7 f_1(4) \} = 0.64$$

$$f_2(8) = \max \{ 0.9 f_1(2), 0.8 f_1(3), 0.7 f_1(5) \} = 0.72$$

$$f_2(9) = \max \{ 0.9 f_1(3), 0.8 f_1(4), 0.7 f_1(6) \} = 0.81$$

$$f_2(10) = \max \{ 0.9 f_1(4), 0.8 f_1(5), 0.7 f_1(7) \} = 0.81$$

$$f_3(10) = \max \{ (0.9) f_2(5), (0.7) f_2(6), (0.5) f_2(8) \} = 0.504$$

So, trace back

& Component 3 $\rightarrow$ 3 parallel, Component 2 $\rightarrow$ 1 parallel, Component 1 $\rightarrow$ 2 parallel.

①4) a) $280 \cdot (1,09)^5 + 210 \cdot (1,09)^4 + 140 \cdot (1,09)^3 + 140 \cdot (1,09)^2 + 210 \cdot (1,09) + 280 =$
 $\times \left((1,09)^4 + (1,09)^3 + (1,09)^2 + (1,09)^1 + (1,09)^0 \right)$

$= 430.815 + 296.432 + 181.304 + 166.334 + 228.9 + 280 = 1583.785$

$\Rightarrow X \cdot \left(\frac{(1,09)^5 - 1}{1,09 - 1} \right) = 1583.785 \Rightarrow X = 264.639$

b) Start of 1st year $2000 \cdot 1,1 + 2000 = 4200$

Start of 2nd year $4200 \cdot 1,1 - 2850 = 1770$

Start of 3rd year $1770 \cdot 1,1 + 3500 = 5447$

Start of 4th year $5447 \cdot 1,1 = 5991,7$