

IE 400: Principles of Engineering Management

In-Class Example Solutions

October 24, 2020

1 Work Scheduling (page:13)

The mathematical formulation of the problem is given as follows:

Decision Variables

x_i : number of workers who are NOT working on day i , $i = 1, \dots, 7$ (1: Mon, 7: Sun)

Objective is to minimize total number of of workers:

$$\min \quad x_1 + x_2 + x_3 + x_4 + x_5 + x_6 + x_7 \quad (1)$$

$$s.t \quad x_2 + x_3 + x_4 + x_5 + x_6 + x_7 \geq 8 \quad (2)$$

$$x_1 + x_3 + x_4 + x_5 + x_6 + x_7 \geq 6 \quad (3)$$

$$x_1 + x_2 + x_4 + x_5 + x_6 + x_7 \geq 7 \quad (4)$$

$$x_1 + x_2 + x_3 + x_5 + x_6 + x_7 \geq 6 \quad (5)$$

$$x_1 + x_2 + x_3 + x_4 + x_6 + x_7 \geq 9 \quad (6)$$

$$x_1 + x_2 + x_3 + x_4 + x_5 + x_7 \geq 11 \quad (7)$$

$$x_1 + x_2 + x_3 + x_4 + x_5 + x_6 \geq 9 \quad (8)$$

$$x_1, x_2, x_3, x_4, x_5, x_6, x_7 \geq 0 \quad (9)$$

$$x_i = \text{integer} \quad \forall i = 1, \dots, 7 \quad (10)$$

2 The 0-1 Knapsack Problem (page:19)

The mathematical formulation of the problem is given as follows:

Decision Variables

$$x_i = \begin{cases} 1, & \text{if item } i \text{ is packed into the suitcase} \\ 0, & \text{otherwise} \end{cases}$$

Objective is to maximize total utility:

$$\max \quad \sum_{i=1}^n c_i x_i \quad (11)$$

$$s.t \quad \sum_{i=1}^n a_i x_i \leq b \quad (12)$$

$$x_i \in \{0, 1\} \quad \forall i = 1, \dots, n \quad (13)$$

3 Workforce Scheduling (page:22)

The mathematical formulation of the problem is given as follows:

Decision Variables

w_t : number of workers working during time period t , $t = 0, \dots, 4$

f_t : number of workers fired at the beginning of time period t , $t = 1, \dots, 4$

h_t : number of workers hired at the beginning of time period t , $t = 1, \dots, 4$

i_t : number of pairs of shoes in the inventory at the end of time period t ,
 $t = 0, \dots, 4$

p_t : number of pairs of shoes produced during time period t , $t = 1, \dots, 4$

Objective is to find a work schedule that meets the total demand at a minimum cost:

$$\min \quad \sum_{t=1}^4 5p_t + 30i_t + 1500w_t + 2000f_t + 1600h_t \quad (14)$$

$$s.t \quad w_0 \geq 3 \quad (15)$$

$$i_0 = 50 \quad (16)$$

$$i_{t-1} + p_t = d_t + i_t \quad \forall t = 1, \dots, 4 \quad (17)$$

$$w_t = w_{t-1} + h_t - f_t \quad \forall t = 1, \dots, 4 \quad (18)$$

$$p_t - 40w_t \leq 0 \quad \forall t = 1, \dots, 4 \quad (19)$$

$$w_t, f_t, h_t, i_t, p_t \geq 0 \quad \forall t = 1, \dots, 4 \quad (20)$$

$$w_t, f_t, h_t, i_t, p_t = integer \quad \forall t = 1, \dots, 4 \quad (21)$$

4 Modeling Fixed Costs (page:68)

The mathematical formulation of the problem is given as follows:

Decision Variables

$$y_i = \begin{cases} 1, & \text{if a potential warehouse } i \text{ is opened } i=1,\dots,m \\ 0, & \text{otherwise} \end{cases}$$

x_{ij} = Number of units supply from warehouse i to client j , $i=1,\dots,m$, $j=1,\dots,n$

Objective is order to minimize total fixed costs and transportation costs while meeting the demand and capacity constraints:

$$\min \quad \sum_{i=1}^m f_i y_i + \sum_{i=1}^m \sum_{j=1}^n c_{ij} x_{ij} \quad (22)$$

$$s.t \quad \sum_{i=1}^m x_{ij} \geq d_j \quad \forall j = 1, \dots, n \quad (23)$$

$$\sum_{j=1}^n x_{ij} \leq a_i y_i \quad \forall i = 1, \dots, m \quad (24)$$

$$y_i \in \{0, 1\} \quad \forall i = 1, \dots, m \quad (25)$$

$$x_{ij} \geq 0 \quad \forall i = 1, \dots, m, \forall j = 1, \dots, n \quad (26)$$

5 The Set Covering Problem (page:70)

The mathematical formulation of the problem is given as follows:

Decision Variables

$$x_i = \begin{cases} 1, & \text{if a hospital is built in potential location } i, i=1,2,3,4 \\ 0, & \text{otherwise} \end{cases}$$

Objective is order to minimize number of hospitals that should be opened:

$$\min \quad x_1 + x_2 + x_3 + x_4 \quad (27)$$

$$s.t \quad x_1 + x_3 \geq 1 \quad (28)$$

$$x_2 + x_3 \geq 1 \quad (29)$$

$$x_2 + x_4 \geq 1 \quad (30)$$

$$x_1 + x_4 \geq 1 \quad (31)$$

$$x_3 + x_4 \geq 1 \quad (32)$$

$$x_i \in \{0, 1\} \quad \forall i = 1, \dots, 4 \quad (33)$$

6 Location Problems (Cover, Center, Median) (page:71-77)

Ankara municipality is providing service to 6 regions. The travel times (in minutes) between these regions, say t_{ij} values are:

	R1	R2	R3	R4	R5	R6
R1	0	10	20	30	30	20
R2	10	0	25	35	20	10
R3	20	25	0	15	30	20
R4	30	35	15	0	15	25
R5	30	20	30	15	0	14
R6	20	10	20	25	14	0

- (a) **Cover:** The municipality wants to build fire stations in some of these regions. Find the minimum number of fire stations required if the municipality wants to ensure that there is a fire station within 15 minute drive time from each region.

Decision Variables: The coverage drive time $T^c = 15$ is the maximum drive time between a fire station and a region that it can serve. Furthermore, in this formulation we do not need to explicitly use the travel times. Instead, we can use the following indicator variable.

$$a_{ij} = \begin{cases} 1, & \text{if } t_{ij} < 15 \\ 0, & \text{otherwise} \end{cases} \quad (34)$$

Let's define the decision variable.

$$x_j = \begin{cases} 1, & \text{if fire station located at } j \\ 0, & \text{otherwise} \end{cases} \quad (35)$$

IP Model:

$$\min \quad \sum_{j=1}^6 x_j \quad (36)$$

$$s.t \quad \sum_{j=1}^6 a_{ij} x_j \geq 1 \quad \forall i = 1, \dots, 6 \quad (37)$$

$$x_j \in \{0, 1\} \quad \forall j = 1, \dots, 6 \quad (38)$$

- (b) **Median:** The municipality now wants to build 2 post offices. There is going to be a back and forth trip each day from each post office to each one of the regions assigned to this post office. Formulate a model which will minimize the total trip time.

Decision Variables:

$$y_{ij} = \begin{cases} 1, & \text{if region } i \text{ assigned to post office located at } j \\ 0, & \text{otherwise} \end{cases} \quad (39)$$

$$x_j = \begin{cases} 1, & \text{if post office located at } j \\ 0, & \text{otherwise} \end{cases} \quad (40)$$

IP Model:

$$\min \quad \sum_{j=1}^6 \sum_{i=1}^6 2t_{ij} y_{ij} \quad (41)$$

$$s.t \quad \sum_{j=1}^6 y_{ij} = 1 \quad \forall i = 1, \dots, 6 \quad (42)$$

$$y_{ij} \leq x_j \quad \forall i, j = 1, \dots, 6 \quad (43)$$

$$\sum_{j=1}^6 x_j = 2 \quad (44)$$

$$x_j \in \{0, 1\} \quad \forall j = 1, \dots, 6 \quad (45)$$

$$y_{ij} \in \{0, 1\} \quad \forall i, j = 1, \dots, 6 \quad (46)$$

- (c) What if additional to part (b), you have the restriction that a post office cannot provide service to more than 4 regions including itself?

Add the following constraint to above model:

$$\sum_{i=1}^6 y_{ij} \leq 4 \quad \forall j = 1, \dots, 6$$

- (d) **Center:** The municipality wants to build 2 fire stations so that the time to reach the farthest region is kept at minimum.

Decision Variables:

$$y_{ij} = \begin{cases} 1, & \text{if region } i \text{ assigned to post office located at } j \\ 0, & \text{otherwise} \end{cases} \quad (47)$$

$$x_j = \begin{cases} 1, & \text{if post office located at } j \\ 0, & \text{otherwise} \end{cases} \quad (48)$$

Let's define an additional decision variable T that represents the maximum travel time between any region and its servicing fire station.

IP Model:

$$\min \quad T \quad (49)$$

$$s.t \quad \sum_{j=1}^6 y_{ij} = 1 \quad \forall i = 1, \dots, 6 \quad (50)$$

$$y_{ij} \leq x_j \quad \forall i, j = 1, \dots, 6 \quad (51)$$

$$\sum_{j=1}^6 x_j = 2 \quad (52)$$

$$\sum_{j=1}^6 t_{ij} y_{ij} \leq T \quad (53)$$

$$x_j \in \{0, 1\} \quad \forall j = 1, \dots, 6 \quad (54)$$

$$y_{ij} \in \{0, 1\} \quad \forall i, j = 1, \dots, 6 \quad (55)$$

- (e) Consider the median problem in part (b). Assume that in addition to the back and forth trips between the post offices and the assigned regions, there will be a trip between the two post offices. The municipality wants to locate the offices in such a way that the total trip time is minimized.

Decision Variables:

Let's define an additional decision variable

$$z_{ij} = \begin{cases} 1, & \text{if post office located at both } i \text{ and } j \text{ } (j \neq i) \\ 0, & \text{otherwise} \end{cases} \quad (56)$$

IP Model:

$$\min \sum_{j=1}^6 \sum_{i=1}^6 2t_{ij}y_{ij} + \sum_{j=1, j \neq i}^6 \sum_{i=1}^6 2t_{ij}z_{ij} \quad (57)$$

$$s.t \quad \sum_{j=1}^6 y_{ij} = 1 \quad \forall i = 1, \dots, 6 \quad (58)$$

$$y_{ij} \leq x_j \quad \forall i, j = 1, \dots, 6 \quad (59)$$

$$\sum_{j=1}^6 x_j = 2 \quad (60)$$

$$z_{ij} \geq x_i + x_j - 1 \quad \forall i, j = 1, \dots, 6, j \neq i \quad (61)$$

$$x_j \in \{0, 1\} \quad \forall j = 1, \dots, 6 \quad (62)$$

$$y_{ij} \in \{0, 1\} \quad \forall i, j = 1, \dots, 6 \quad (63)$$

7 Blending problem revisited with binary variables (page:83-85)

A company produces animal feed mixture. The mix contains 2 active ingredients and a filler material. 1 kg of feed mix must contain a minimum quantity of each of the following 4 nutrients:

Nutrients	A	B	C	D
Required amount (in grams) in 1 kg of mix	90	50	20	2

The ingredients have the following nutrient values and costs:

	A	B	C	D	Cost/kg
Ingredient1 (gram/kg)	100	80	40	10	\$40
Ingredient2 (gram/kg)	200	150	20	N/A	\$60

- (a) What should be the amount of active ingredients and filler material in 1 kg of feed mix in order to satisfy the nutrient requirements at minimum total cost?

Decision Variables:

x_1 = amount (kg) of ingredient 1 in one kg of feed mix

x_2 = amount (kg) of ingredient 2 in one kg of feed mix

x_3 = amount (kg) of filler in one kg of feed mix

IP Model:

$$\begin{array}{ll}
\min & 40x_1 + 60x_2 & (64) \\
s.t & x_1 + x_2 + x_3 = 1 & (65) \\
& 100x_1 + 200x_2 \geq 90 & (66) \\
& 80x_1 + 150x_2 \geq 50 & (67) \\
& 40x_1 + 20x_2 \geq 20 & (68) \\
& 10x_1 \geq 2 & (69) \\
& x_1, x_2, x_3 \geq 0 & (70)
\end{array}$$

(b) Suppose now that we have the following additional limitations

- If we use any of ingredient 2, we incur a fixed cost of \$15 (in addition to the \$60 unit cost).
- It is enough to only satisfy 3 of the nutrition requirements rather than 4.

Decision Variables (Additional to those for Part a):

$$y = \begin{cases} 1, & \text{if } x_2 \geq 0 \\ 0, & \text{otherwise} \end{cases} \quad (71)$$

$$z_i = \begin{cases} 1, & \text{if nutrient constraint } i \text{ is satisfied} \\ 0, & \text{otherwise} \end{cases} \quad (72)$$

IP Model:

$$\min \quad 40x_1 + 60x_2 + 15y \quad (73)$$

$$s.t \quad x_1 + x_2 + x_3 = 1 \quad (74)$$

$$100x_1 + 200x_2 \geq 90z_1 \quad (75)$$

$$80x_1 + 150x_2 \geq 50z_2 \quad (76)$$

$$40x_1 + 20x_2 \geq 20z_3 \quad (77)$$

$$10x_1 \geq 2z_4 \quad (78)$$

$$x_2 \leq y \quad (79)$$

$$z_1 + z_2 + z_3 + z_4 \geq 3 \quad (80)$$

$$x_1, x_2, x_3 \geq 0 \quad (81)$$

$$y \in \{0, 1\} \quad (82)$$

$$z_i \in \{0, 1\} \quad \forall i = 1, \dots, 4 \quad (83)$$